

# A STUDY ON REVIEW AND EVALUATION OF APPROXIMATE REVERSE CARRY PROPAGATE ADDER

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## **ABSTRACT:**

*An Adder play very important position into several digital signals processing application & can efficiently (RCPA) are presents into this papers in which is a Reverse carry propagate adders ( RCPA) structures, the carry signal propagate at counter flow manners for the most significantly bits as the least significantly bits. Therefore, the carry input signals have high meaning to the output carry. The technique of the carry propagations leads as to high stability to the incidence into delay variation. At this time is a papers as proposed in three several design in the reverse carry propagate full adder (RCPFA) cells within many delays, energy, & accuracy configuration. It is also uses in a number of the hybrid structures into the n-bit adder designed have been least half to the adder designs are implement for this RCPA & most significantly half would be add through some accurately higher speeds adder such as Kogge stone adders in which is enhanced in the adder speeds. The Hybrid adder realizes as utilizes in this structure are study & compare in those within convention approximate adder using Xilinx ISE 14.7 for Verilog HDL coding.*

**Keywords:** *Approximate adders, digital signal processing (DSP), Reverse carry propagate adders (RCPA) etc.*

## **I. INTRODUCTION**

The electricity utilization decrease and speed development are the important thing objectives within the shape of automatic managing devices, mainly the compact frameworks. Regularly, an expansion within the pace is finished at the price of greater power utilization for careful handling devices. One of the methods to address enhance every the energy and tempo is to forfeit the calculation precision this device, this is surmised processing, might be carried out for the programs in which a few errors is probably continued. They incorporate those where virtual sign Processing (DSP) are carried out at the human revel in-related signs and signs and signs and signs and symptoms thinking about the reality that human perceptual capacities are confined, a huge a part

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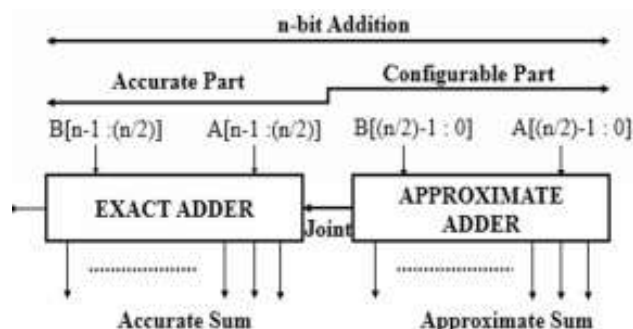
of the sports, the surmised registering is probably conjured for custom DSP squares which coping with the ones signs and symptoms and signs and symptoms.

Adder blocks, which might be the principle segments in range juggling gadgets of DSP frameworks, are managing hungry and frequently shape problem place areas at the pass on. the ones actualities were the inspirations for information this issue the use of the surmised registering method in advance seems into on difficult adders have decided modern strategies of specializing in blunder weight and blunder hazard decreases. The primary method is based upon on a aggregate form viper where in precise elements, accurate MSBs, and tough least huge bits (LSBs) are used. The blunder shows up in the supply contribution of the best most sizable bit (MSB) detail and the summation within the LSB difficulty. This constrains the blunder load to the heaviness of the bring contribution of the MSB element. due to the fact frequently the fantastic majority of the wearing activities get up inside the LSB detail, manipulate decreases over 70% might be completed the usage of the crossover adder method inside the 2nd technique, unadulterated inexact adder structures are implemented. For those adders, reducing the mistake threat of the summation truly as diminishing the electricity and postponement are the vital thing plan standards. They will likewise be joined thru a mistake exchange unit which has time, power, and territory overheads.

We centre at some stage in the adders in which the usage of the tough switch opposite deliver Propagate whole Adder (RCPFA) is proposed. The predicted viper proliferates the statistics supply in a counter-circulate way, i.e., from the better important piece to supply down noteworthy piece to border the supply yield. in this kind of adder, that's referred to as turn round carry opposite deliver Propagate Adder (RCPA), the proliferation is carried out through using offering a conjecture flag going approximately as a yield flag. Inferable from the turnaround engendering, the heaviness of the supply diminishes as it spreads. This shape of adder improves the deferral and energy contrasted with the ones of the reducing issue surmised adders. Moreover, this adder kind is a whole lot less defenceless inside the course of the remove variety when contrasted with the commonplace ones special widely identified of the proposed RCPFA are contrasted with the ones of the inexact FAs. Manages exploring the plan parameters of the proposed FAs and the adequacy of their utilization in a blunder bendy software application.

## II. EXSITING TECHNIQUE

Among all the exact adder systems, a Ripple deliver adder (RCA) has the worst strength and region utilizations. It, however, suffers from a big do away with. to enhance the rate and energy performance of this adder, a few in advance works have sacrificed the accuracy. An approximate RCA shape which modified into called mistakes tolerant adder form (EAS) has been provided. The shape of EAS is shown in Fig. 1. On this form, the input operands are divided into components known as genuine computation element and inexact computation thing. inside the correct element is the MS element, the conventional FAs with a 0 deliver input for the complete element are used even as the inexact thing is the LS element which includes a deliver-free addition detail.



**Fig.1 Error tolerant adder structure**

Designing an adder with supply propagation will take greater do away with so if we avoid the supply propagation while designing an adder will outcomes the better adder common ordinary performance so we pass for contrary deliver propagate adder in which the supply is propagated in contrary route. The conventional FA that is the crucial issue building block of the deliver propagates adders has 3 inputs with the identical weight. Moreover, it has outputs for a summation give up cease end result with the same weight as that of the inputs and a deliver output with times the burden. The deliver propagation postpone (i+1) is the most vital timing parameter in an FA due to the reality that it determines the put off of the crucial path of multi bit adders (and multipliers). A small short-put off violation may also additionally result in a massive quantity of errors due to the fact that the error takes vicinity on the MSBs of the summation this is the stop end result of the technology and propagation of the carry input of the MSBs via small sizable bit FAs primarily based mostly on this reasoning, if the order of the supply propagation is reversed, one can also moreover assume that the amount of errors due to the timing violation decreases.

TABLE I

ERROR EVALUATION METRICS OF DIFFERENT APPROXIMATE ADDERS

	ER %	MED	MRED	Max ED	$\mu$	$\sigma$
RCPFA I	75.95	18.20	0.4439	128	-0.33	31.98
RCPFA II	<b>65.99</b>	<b>18.12</b>	0.4763	<b>127</b>	18.12	<b>26.57</b>
RCPFA III	80.08	18.79	<b>0.4392</b>	128	-18.79	26.58
ETA I	89.99	51.18	0.8384	255	51.18	57.11
AMA I	85.91	34.57	0.9429	255	<b>0</b>	60.36
AMA II	89.99	59.69	1.1235	255	-2.00	85.31
AMA III	97.56	71.31	1.5805	255	-1.00	94.03
AMA IV	97.56	66.29	3.8755	255	31.75	82.62
TGA I	89.99	44.79	2.5951	170	-0.25	64.00
TGA II	85.91	32.25	0.7780	170	-32.25	45.25
AXA-I	96.66	255.5	12.0102	510	255.5	128
LOA	89.64	47.69	0.7910	128	-0.249	63.88

It, however, suffers from a large put off to decorate the speed and power overall performance of this adder; a few earlier works have sacrificed the accuracy. In [5], an approximate RCA form which have become referred to as errors tolerant adder kind I (ETA I) come to be supplied. The shape of ETA I is confirmed in Fig. 1 in this form, the input operands are divided into particular MS and inexact LS additives inside the real MS

issue, the conventional FAs with a zero deliver input for the complete detail are used whilst the inexact LS detail includes a deliver-loose addition element (which incorporates XORs) and a control block. The manage block devices all the give up result bits to “1” from the wonderful bit characteristic on the inexact problem in which each of the corresponding bits of the inputs are “1” (component B) to the LSBs of the inputs. Moreover, the result bits from the element B to the turning into a member of component are generated by means of the usage of using the usage of the bring-free addition. In [6], the overall adder of the LS a part of the adder has been modified thru OR gates vital to smaller put off, strength intake, and region. Moreover, an AND gate has been employed to generate the enter bring of the MS element. In [7], 5 approximate replicate adder (AMA) structures having smaller fashion of transistors in contrast to that of the conventional adder were proposed these designs had been based totally on simplifying the inner shape (getting rid of the transistors) of the replicate adder essential to smaller area and strength intake similarly to higher tempo. The truth tables of AMA-I to AMA-IV are depicted in table I. In AMA-V, the sum and deliver outputs are right away related to the inputs fending off the usage of any complete adder on the equal time as this form is rapid and ultralow strength, its accuracy is very low.

### III. PROPOSED TECHNIQUE

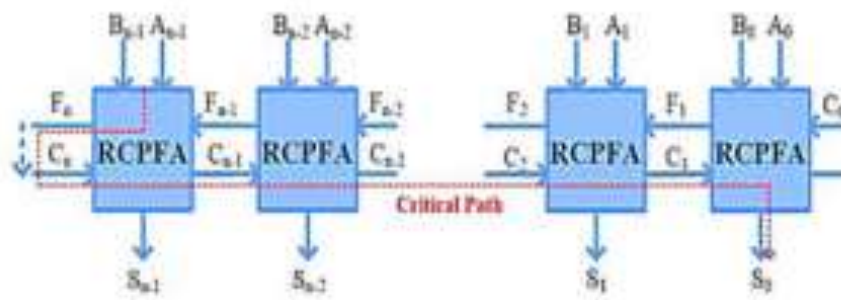
Reverse Carry Propagate Full-Adder Cell Each exact FA generates its carry output and sum signals using

$$2C_{i+1} + S_i = A_i + B_i + C_i \dots \dots \dots (1)$$

Where (A<sub>i</sub>) is the bit of the input (i), (i+ 1)is the carry input (output), and S<sub>i</sub> is the bit of the sum. Based on this equation, the output signals in the bit position depends on the bits of the inputs A and B and the carry output of the previous position (B<sub>i</sub>). By moving the term (i+1)to the left (right) side of the equation, one may write

$$S_i - C_i = A_i + B_i - 2C_{i+1} \dots \dots \dots (2)$$

Thinking about (2), one might imagine of a complete adder as a shape which operation depends at the deliver output of the (A<sub>i</sub> + 1) nth bit role (B<sub>i</sub> + 1) and its enter operand bits. For this shape, the outputs are the sum and the convey signs and symptoms with the identical weights based on the above talk, we endorse a own family of complete adders for the RCPFA established in Fig. 2. As proven in Fig. 2, those full adders have four inputs and 3 outputs. The inputs are the enter operands (A<sub>i</sub> and B<sub>i</sub>), the deliver output of the following bit feature (C<sub>i+1</sub>), and a forecast sign (F<sub>i</sub>). The RCPFA determines the summation surrender end result (S<sub>i</sub>), deliver (C<sub>i</sub>), and the forecast sign (F<sub>i+1</sub>) as its output indicators. As mentioned earlier as, the benefit of the RCPA is that the price of the mistake is inside the route of lower within the bit importance because of this the cumulative effect of the mistake (e.g., because of the take away variant) at some stage in the convey propagation is decrease for bits with better significances

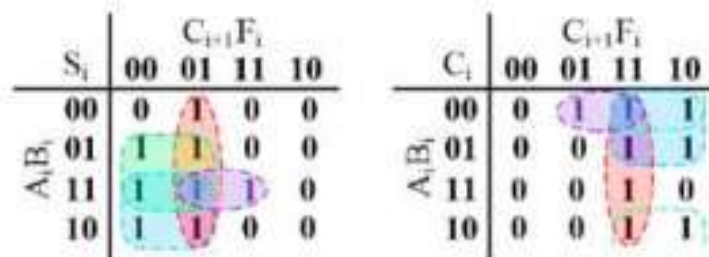


**Fig.2 The n-bit RCPA**

To determine a structure for RCPFA, the Karnaugh maps of the summation result ( $S_i$ ) and carry ( $C_i$ ) were drawn based on (2) and considering the forecast signal as an input (Fig. 2) The Boolean relations between inputs for generating  $S_i$  and  $C_i$  are obtained as

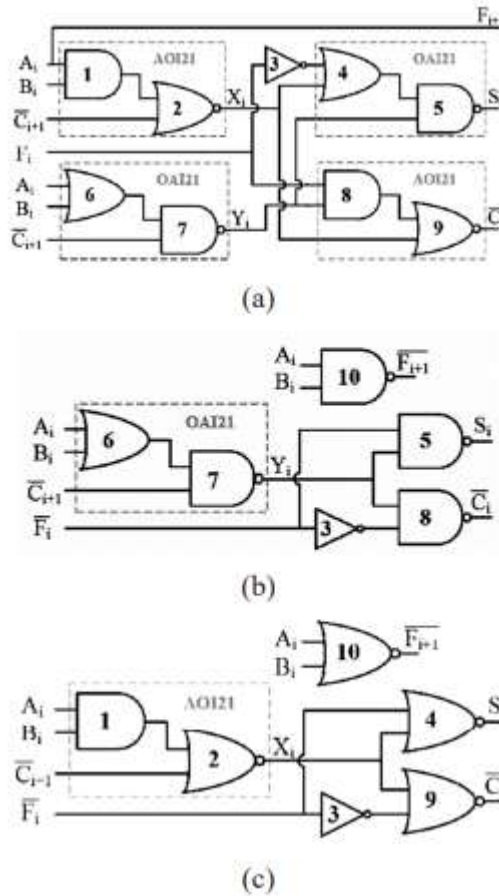
$$S_i = \overline{C_{i+1}}F_i + \overline{C_{i+1}}A_i + \overline{C_{i+1}}B_i + A_i B_i F_i \dots \dots (3)$$

$$C_i = C_{i+1}F_i + C_{i+1}\overline{A_i} + C_{i+1}\overline{B_i} + \overline{A_i}\overline{B_i}F_i \dots \dots (4)$$



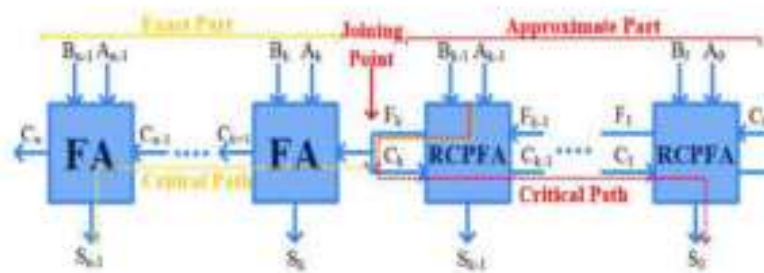
**Fig.3 Karnaugh maps for signals  $S_i$  and  $C_i$  of the general form of RCPFA.**

The primarily based on our requirement we layout 3 forms of opposite convey propagate adders the designs of these adders is given underneath. by using the use of those proposed adders we are able to enforce 3 sorts of opposite propagate adders and are used in the vicinity of approximate adders that's proven on figure 3.



**Fig.4 Internal structures of the (a) RCPFA-D1, (b) RCPFA-D2, (c) RCPFA-D3**

The block diagram of an n-bit adder which is designed with the help of reverse carry propagate adder is shown in below figure As mentioned before, the weight of the carry decreases as the carry propagates in a counter-flow manner.



**Fig.5 Architecture of an n-bit adder with RCPA**

This belonging permits having a lot less vulnerability to the put off variation (because of manner and supply voltage versions) impact for this adder compared to awesome proposed approximate FAs in that is especially excessive first-rate within the case of hybrid adders with huge sizes for the approximate issue which determines the crucial route of the adder. The proposed RCPFAs may be utilized in hybrid adders whose modern day n-bit structure based totally on the RCPFAs is depicted in Fig. 5. Obviously, the layout parameters of the adder rely on the width of the approximate thing.

In present framework, thinking about there is numerous dangers the flip round conveys proliferate adder is step by step gifted. To shape a fast territory inexperienced low power low power half and 1/2 of 64-bit adder for future DSP software program. Increment the speed of adder thru making use of low doors and transistor (22n CMOS innovation) lower the electricity utilization thru lessening the power deliver to 0.8 V. Adder rectangular is the center of ALU, multiplier and processors increasing the price of snake will construct the price of ALU, multiplier and processors. The power utilization lower and pace development may be finished. It improves the deferral and energy. It is a lot less helpless against postponement contrasted with the convectional ones. The proposed technique consists of ease shape. The deferral and mistakes can be redressed. The power usage lower and tempo development can be finished. RCPA improves the postponement and electricity than inexact adders. It is a first rate deal less helpless in opposition to deferral contrasted with the convectional ones. These effects in growing the fee of the interest of DSP. A massive a part of the squares is selected thru ALU and multiplier. Increasing the speed of adder will gather the velocity of ALU and multiplier. Adder block is the middle of ALU and multiplier.

#### IV. RESULTS AND DISCUSSION

By implementing the adder with two sub adders like Kogge stone adder and reverse carry propagate adder we can have better in terms of delay. Here in this paper we present 16-bit and 32-bit adders with 3 different designs of reverse carry propagate adders. The results are presented in the below table

Table 2 Comparison between Existing and proposed Adders

Adder Type	Existing		Proposed	
	Area	Delay	Area	Delay
16bit_d1	45	13.341	62	13.003
16bit_d2	34	13.322	50	10.982

16bit_d3	33	13.308	48	11.296
32bit_d1	95	21.829	148	22.461
32bit_d2	75	21.809	132	15.318
32bit_d3	80	21.765	137	16.296

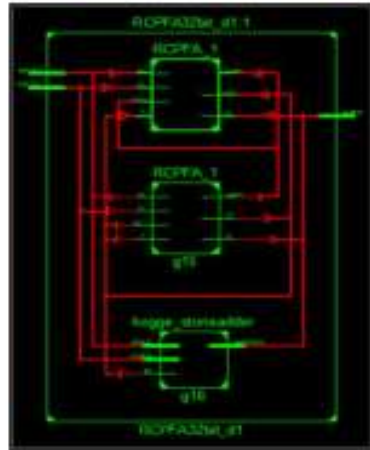


Fig. 6 RTL Schematic of proposed 32 bit adder design 1

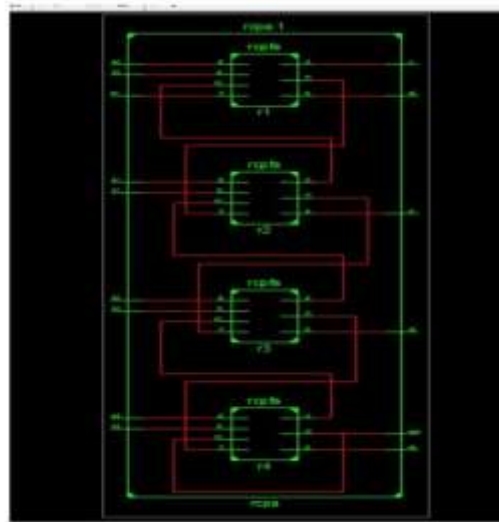


Fig.7. Ripple-Carry adder using RCPFA Cell.

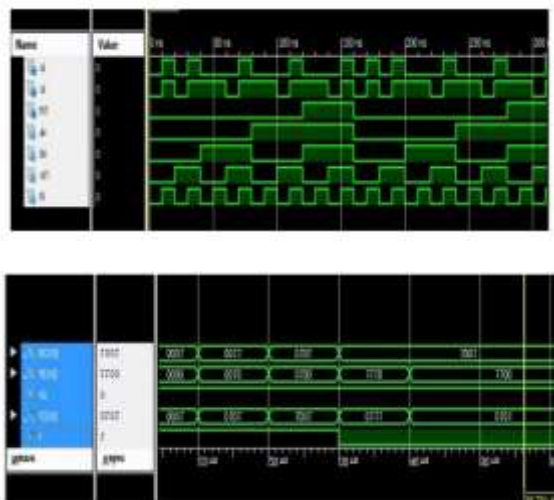


Fig 8. Simulation Result for RCPFA



## V. CONCLUSION

To this paper we've were given designed and accomplished an n-bit adder with a hybrid structure in which we use 2 sorts of adders for implementation of huge adder. We use approximation in lower aspect addition this is least large aspect and correct addition in higher aspect that is most large aspect. The correct adder we used in this assignment is a parallel prefix adder referred to as Kogge stone adder and the approximate adder we used is the RCPA we proposed on this paper through this structure we are capable of get higher delay at the same time as degrading the accuracy and area.

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